

CS180 - Asymptotics

Note Title

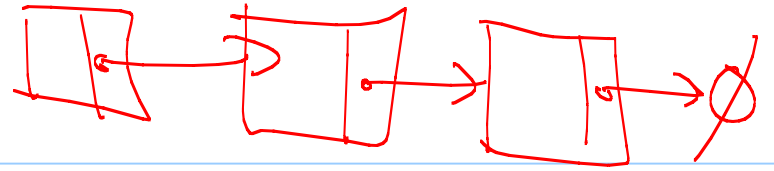
9/22/2011

Announcements

- Lecture tomorrow
- HW due Monday
- Next program will be up Saturday

Ch 3.2 (?)

- head →



Functions

- h A L e

public :

✓ SLinked List ();

~ SLinked List ();

✓ bool empty () const ;

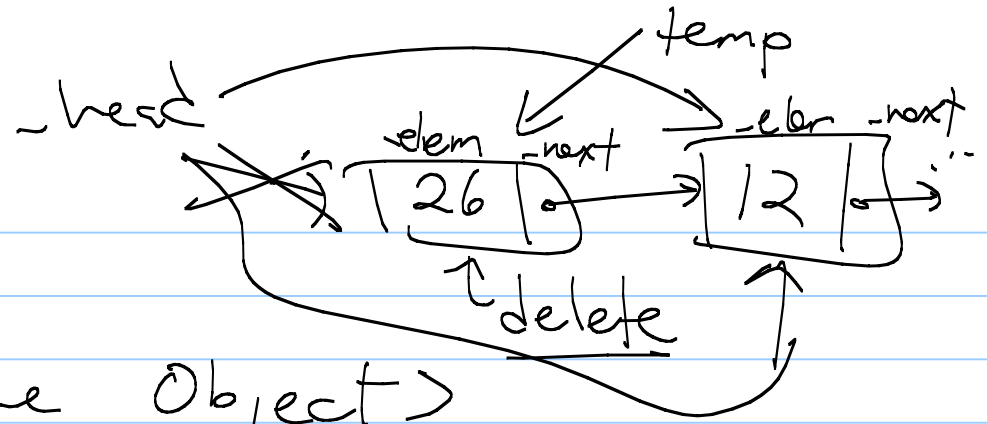
✓ Const Object & front () const ;

void addFront (const Object & e);

✓ void removeFront ();

} ;

remove Front



```
template <typename Object>
```

```
void SLinkedList<Object>::removeFront() {
```

```
    SNode<Object> * temp;
```

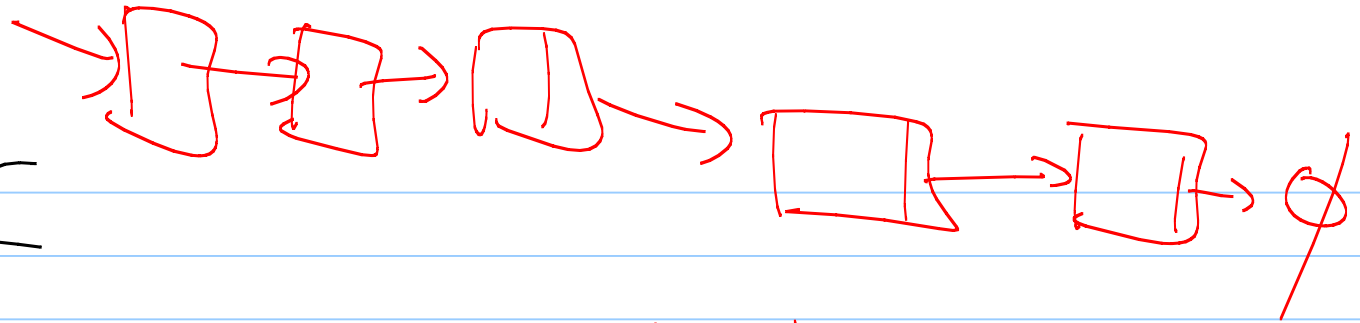
```
    temp = _head;
```

```
    _head = _head -> _next;
```

```
    delete temp;
```

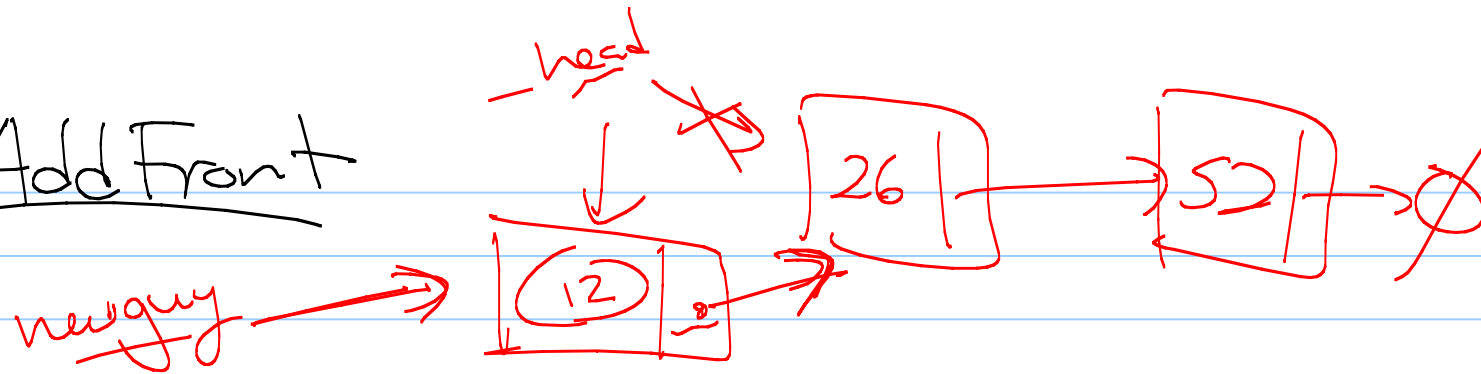
```
}
```

Destructor



```
template < typename Object >
SLinkedList < Object > :: ~SLinkedList() {
    while ( !empty() ) {
        removeFront();
    }
}
```

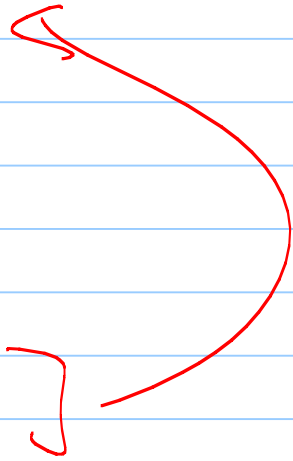
Add Front



```
template < typename Object >
void SLinkedList < Object > :: addFront (const Object (data)) {
    SNode < Object > * newguy = new
        SNode < Object > ;
    newguy -> _elem = data ;
    newguy -> _next = _head ;
    _head = newguy ;
}
```

Algorithm Analysis (Ch. 4)

How do we compare two programs?

- memory usage
 - speed
 - X - interface
 - X - features
 - benchmarks
 - X - cost
- 

Speed

How fast an algorithm runs can be very dependent on variables in the system.

Examples:

- hardware
 - other software on system
 - OS →
 - language
-
- input

Primitive Operations

As a way to compare algorithms in a generic way, we instead count primitive operations.

↳ addition, subtraction, memory access, return, mult & div

In addition, we (generally) only analyze the worst possible running time.

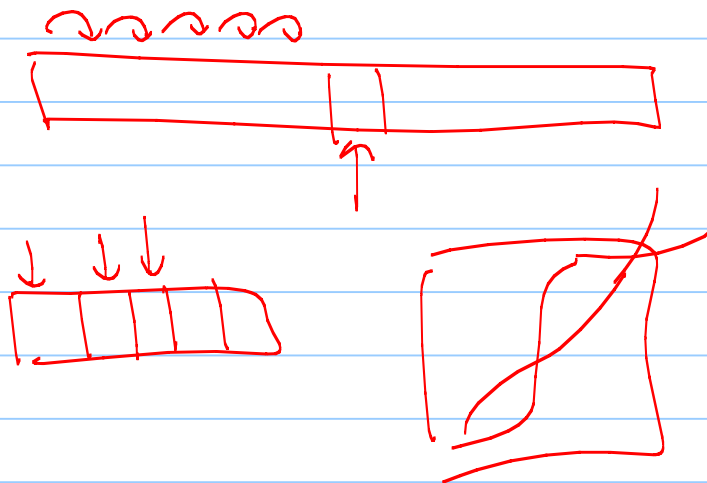
Why? guaranteeing a minimum performance

Comparing

OK, so we have the worst case # of operations - usually a function of n .

How to compare?

Binary search $\sim \log_2 n$
versus
Linear search
↑ checks every element



$4n \log_2 n$ operations
 $2n \log_2 n + 1.265n$

Big-O

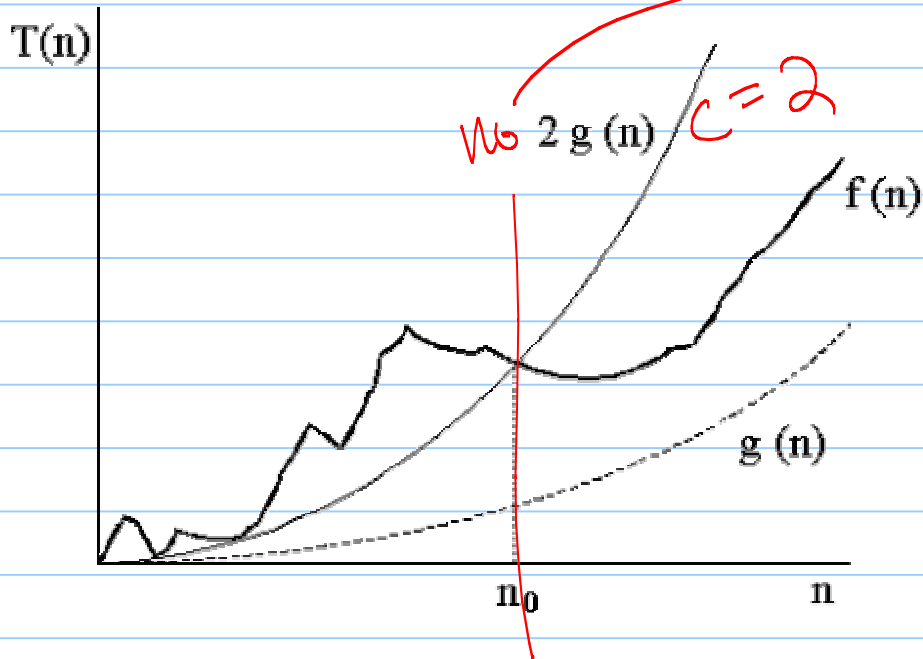
$$f(n) \leq g(n)$$

We say $f(n)$ is $O(g(n))$ if $\forall n > n_0$,

$\exists c > 0$ such that $f(n) \leq c \cdot g(n)$.

for all

there exists



Ex: $5n$ is $O(n^2)$
 $\forall n \geq n_0$, $5n \leq 5n^2 = C$

Ex: $5 \cdot n$ is $O(n)$
 $5n \leq 5n = C$

Ex: $16n^2 + 52$ is $O(n^2)$
 $16n^2 + 52 \leq 16n^2 + 52n^2 = 68n^2 = C$

\Rightarrow $a_n x^n + a_{n-1} x^{n-1} + \dots + a_0$ is $O(x^n)$

Functions we will use

$$\log_2 a + \log_2 b = \log_2(a \cdot b)$$
$$\log_2 X^c = c \log_2 X$$

- ① $O(1)$ - constant time
- ② $O(\log n)$ - logarithmic time
↳ Binary search
- ③ $O(n)$ - linear time
- ④ $O(n \log n)$
- ⑤ $O(n^2)$ - quadratic time
- ⑥ $O(n^3)$ - cubic time
- ⑦ $O(2^n)$ - exponential time

