Network Flow

CSCI 3100

What is Network Flow?

Network Flow is a subsection of graph theory related specifically to situations where something moves from one location to another.

An easily visualized example of network flow is piping system through which a quantity of water must pass.

Concepts in Network Flow

"Sources" are nodes which supply a commodity

"Sinks" are nodes which use up a commodity

An edge's capacity is the maximum amount of flow which can pass through it

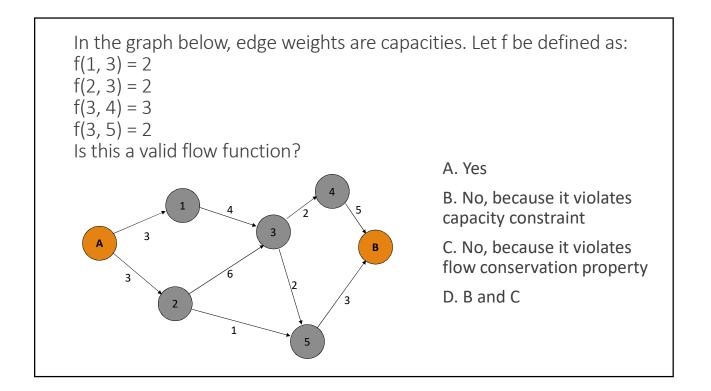
Graphs are usually directed

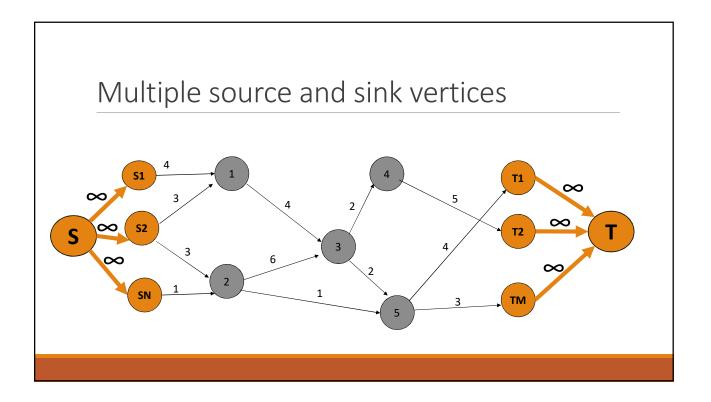
Properties/Constraints

A flow in a graph G = (V, E) is a function f: $V \times V \rightarrow \mathbb{R}$ that satisfies the following properties **Capacity constraint:** for all (u, v) in $V, 0 \le f(u, v) \le c(u, v)$, where c(u, v) is a capacity of edge (u, v)**Flow conservation:** flow into a vertex equals to flow out of that vertex, with the exception of source and sink vertices.

 $\sum_{v \in V} f(v, u) = \sum_{v \in V} f(u, v)$

G is directed In-degree of the source vertex is 0 Out-degree of the sink vertex is 0





Applications of Maximum-Flow

Homework assignment

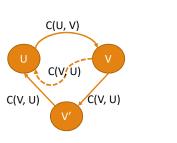
Antiparallel edges

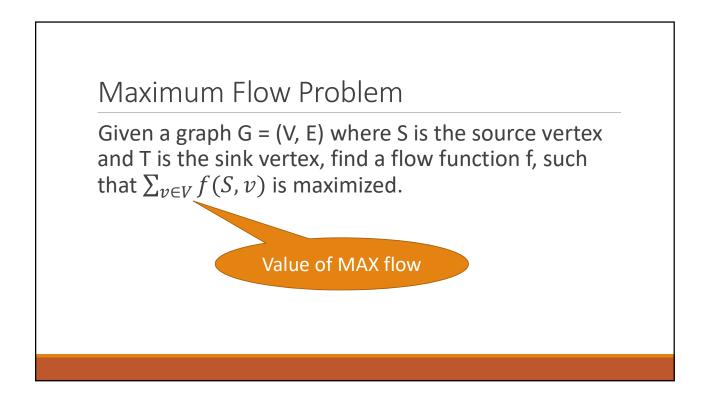
In a flow network, if we have an edge (u, v), we don't want an edge (v, u).

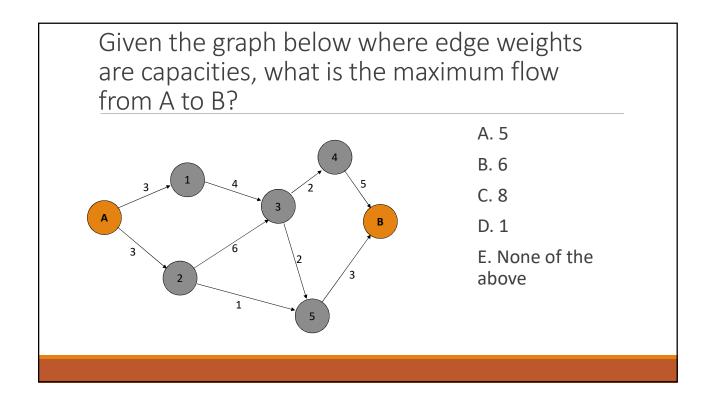
What if we have a network that does not satisfy this constraint?

We can still model this problem as a "flow network":

- $\circ~$ If edges (u, v) and (v, u) are present, introduce a new vertex v' and edges:
- (v, v') and (v', u)
- $\,\circ\,$ The capacity of (v, v') and (v', u) are the same as capacity of (v, u)
- Remove edge (v, u)







The Ford-Fulkerson method

This is a relatively simple way to find maximum flow through a network:

- Find an unsaturated path from the source to the sink
- Add an amount of flow to each edge in that path equal to the smallest capacity in it
- Repeat this process till no more paths can be found
- The total amount of flow added is then maximal